



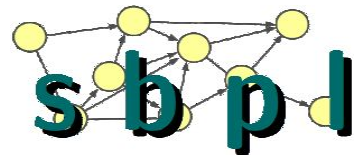
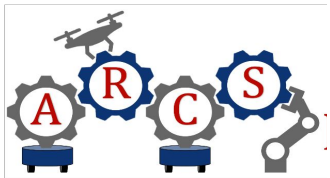
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Windowed MAPF with Completeness Guarantees

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Typical Heuristic Search MAPF Methods

- Many solvers
 - Conflict-Based Search: CBS, ECBS, EECBS, W-EECBS
 - Greedy: PIBT, LaCAM, LaCAM*
 - Others: MAPF-LNS2, BCP

- Majority plan full trajectories

← May not be possible in applications with tight runtime requirements!

Windowed MAPF

- Robot Runners competition had a runtime cutoff of 1 second
- With severe time constraints, **all MAPF works** utilize Windowed planning
 - Only plan a partial collision free path
 - Execute certain # of steps
 - Replan
- Existing works:
 - Windowed Cooperative A*
 - RHCR
 - exPBS
 - Windowed Parallel PIBT-LNS (WPPL)
 - ...

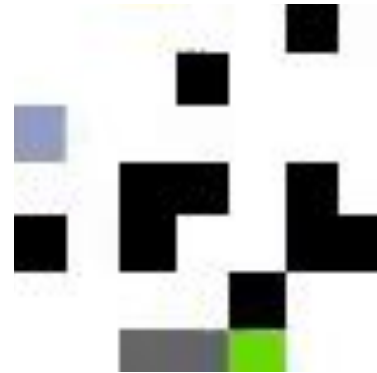
Pros and Cons of Windowed MAPF

Pros:

- Faster planning
- More adaptable to changes
- Enables decentralized planning

Cons:

- Theory: No solution guarantee
- Practice: Myopic reasoning = Deadlock / Livelock



Livelock



Deadlock

Contribution 1: Windowed MAPF with Guarantees!

- We design a framework for windowed MAPF with completeness!
 - **Windowed Complete** MAPF, or WinC-MAPF
- Theory: Completeness guarantees for certain windowed MAPF solvers
- Practice: Overcomes deadlock / livelock

Main ideas:

- Leverage heuristic updates from single-agent real-time search literature
- Leverage agent semi-independence from MAPF literature

Background: Real-Time Heuristic Search (RTHS)

- A problem formulation in single-agent planning where agents can only plan for a finite time (or horizon)
- Their key idea: Maintain completeness by updating the heuristic
 - Suppose at state S , we find S' which minimizes $c(S,S') + h(S')$
 - Update: $h(S) = c(S,S') + h(S')$
 - This heuristic update "penalizes" visited states
 - Heuristic penalty: $h(S) = h_{\text{original}}(S) + h_p(S)$
 - $h_p(S) = c(S,S') + h(S') - h_{\text{original}}(S)$
 - RTHS algorithms prove completeness via the update/penalty

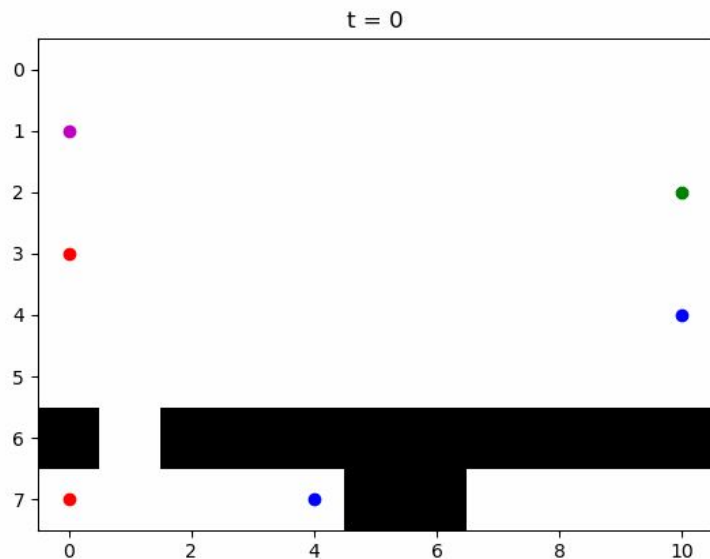
Our key idea: We can reuse this if we view MAPF in the joint configuration

Naive W-MAPF Framework

1. Query "Action Generator" to get limited horizon plan
 - $S \rightarrow S_1 \rightarrow S_2 \rightarrow \dots \rightarrow S_{\text{last}}$
 - AG minimizes $c(S, S_{\text{last}}) + h_{\text{original}}(S_{\text{last}}) + h_p(S_{\text{last}})$
2. Update the heuristic using the standard RTHS update
 - $h_p(S) = c(S, S_{\text{last}}) + h(S_{\text{last}}) - h_{\text{original}}(S)$
3. Execute actions & repeat

Our key idea: View MAPF in the joint configuration

Problem: Heuristic Updates in Joint Space



Has to explore all joint configurations

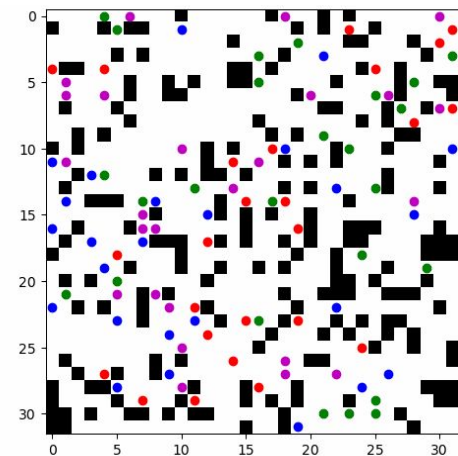
- Involves irrelevant agents!

Solution: Apply Heuristic Updates to Agent Groups

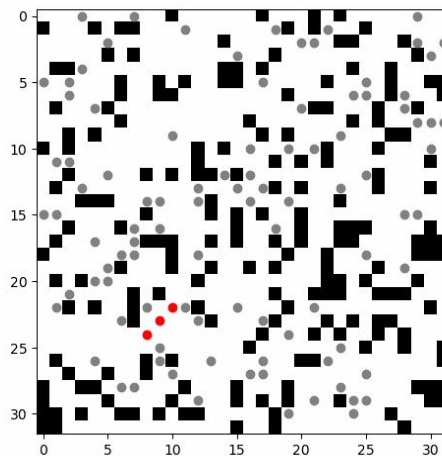
- All performant MAPF methods utilize agent semi-independence
 - I.e. if agents were far apart, they would plan independently
- Examples:
 - CBS: Plans agents independently, then uses conflicts
 - LNS2: Uses prioritized planning where agents plan independently conditioned on others
 - PIBT: Each agent acts greedily, avoids earlier agents
- Solution: Only apply heuristics to "coupled" agent groups
 - Ignore independent agents
 - Note: Paper discusses more formally "disjoint agent groups"

Only apply heuristic penalties to coupled agents

Visualizing Groups in Action



Current configuration



Dynamically Detected
Disjoint Agent Groups

WinC-MAPF Framework

Algorithm 1: Windowed Complete MAPF Framework

```
1: procedure WINDOWED COMPLETE MAPF( $\mathcal{C}^{cur}$ )
2:    $\mathcal{H} \leftarrow \emptyset$  ▷ Heuristic Penalties
3:   while  $\mathcal{C}^{cur} \neq \text{Goal}$  do
4:      $\mathcal{C}', \text{ListOfGroups} = \text{AG}(\mathcal{C}^{cur}, \mathcal{H})$ 
5:     for  $Gr \in \text{ListOfGroups}$  do ▷ For each group
6:        $h_{new} = c(\mathcal{C}_{Gr}^{cur}, \mathcal{C}'_{Gr}) + h(\mathcal{C}'_{Gr})$ 
7:        $penalty = h_{new} - h_{BD}(\mathcal{C}'_{Gr})$ 
8:       if  $penalty > 0$  then
9:          $\mathcal{H}.insert(\mathcal{C}_{Gr}^{cur}, penalty)$ 
10:     $\mathcal{C}^{cur} \leftarrow \mathcal{C}'$  ▷ Move agents
```

Key idea 1 for **completeness**:
Apply heuristic updates on
joint-space!

Key idea 2 for **efficiency**:
Agent Groups!

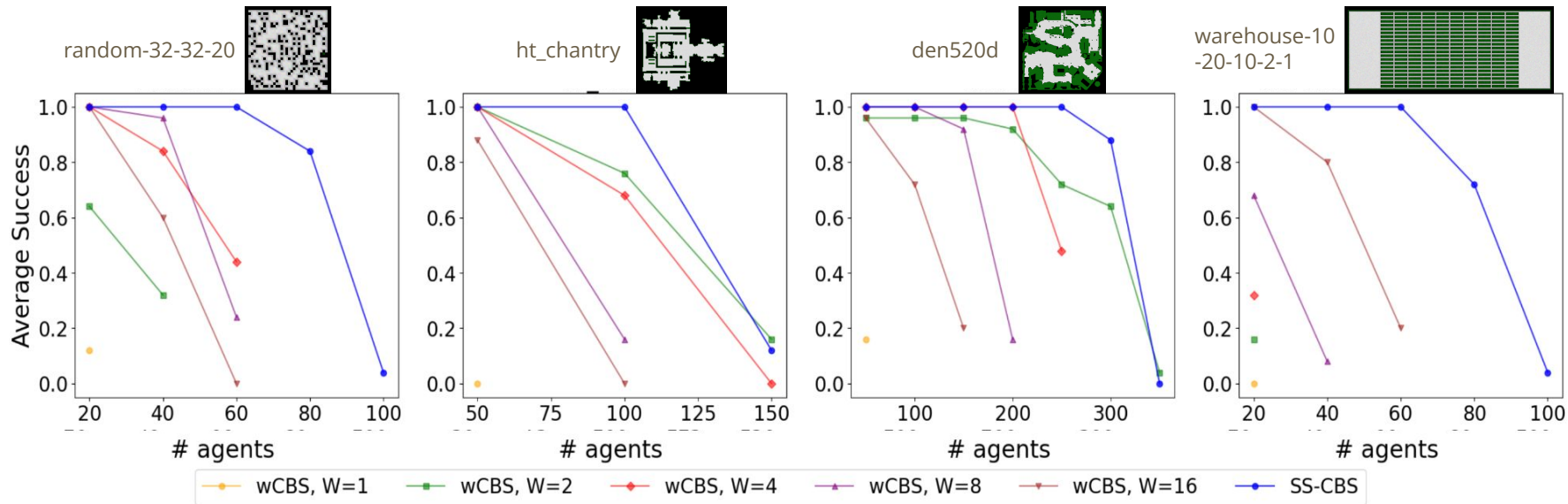
Contribution 2: Single-Step CBS

A concrete instantiation of our framework using CBS

- Detecting Disjoint Agent Groups
- Incorporating Heuristic Penalties

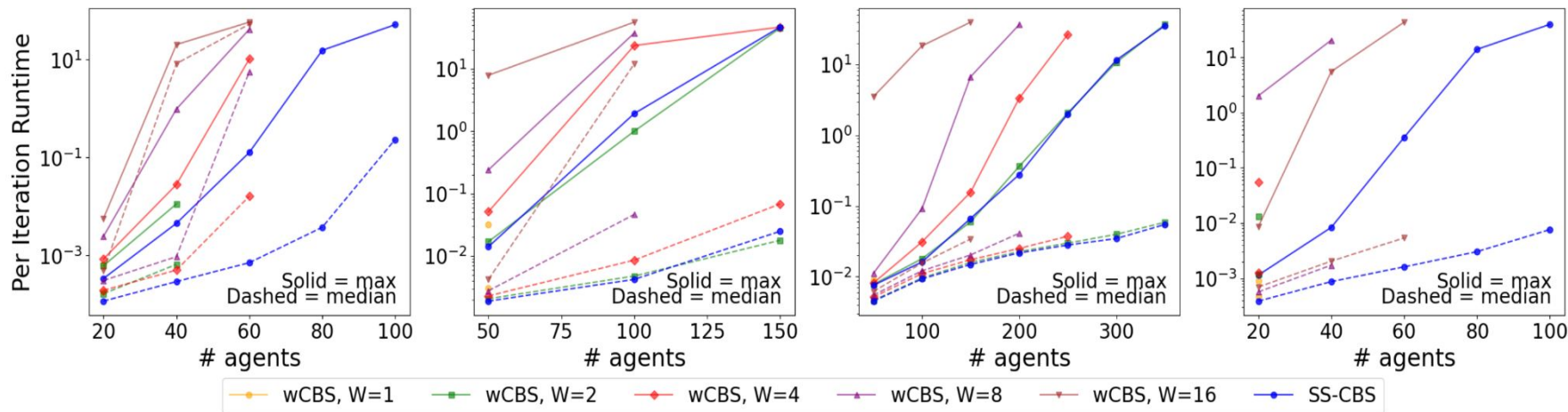
Results

Large Map Success Rate: SS-CBS vs Windowed CBS



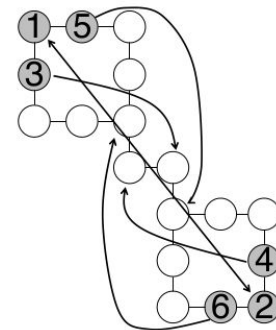
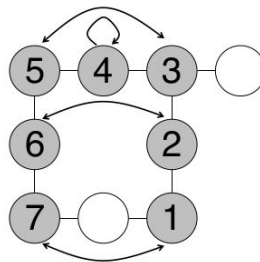
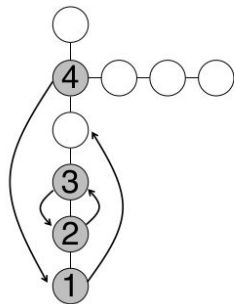
SS-CBS has a higher success rate than existing windowed CBS across multiple window sizes

SS-CBS Runtime



The median iteration is very fast, but certain iterations are bottlenecks

Small Tough Maps



Method	Horizon	Tunnel		Loopchain		Connector	
		3	4	6	7	5	6
wCBS	1,2,4,8,16	-	-	-	-	-	-
CBS+	∞	0.9	-	-	-	-	-
SS-CBS	1	1	1	1	0.95	1	1

+ SS-CBS can solve tough congestion that windowed approaches cannot
 - Note: Solution cost is very poor (100s-1000s)

Takeaways

- Existing Windowed MAPF methods are necessary for runtime performance, but are incomplete
- Our WinC-MAPF Framework enables theoretical completeness
 - And practical improvements
- SS-CBS is an initial WinC-MAPF method
 - Better success rate in small and large instances
- Lots of future work!
 - Multi-step solvers
 - Faster or suboptimal solvers (requires proving completeness)
 - Incorporate single agent heuristic updates as well
 - Lifelong MAPF variants

This work opens up a new line of MAPF research!